Exploration for Free: How Does Reward Heterogeneity Improve Regret in Cooperative Multi-agent Bandits?

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Motivation Example for Action Constrained Multi-Agent Bandits



Agents have access to its nearby arms

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- **1** *M* Agents and *K* Arms (set \mathcal{K})
- 2 Each agent *m* has access to a subset of arms $\mathcal{K}^{(m)} \subseteq \mathcal{K}$
- 3 Overlap $\mathcal{K}^{(m)} \cap \mathcal{K}^{(m')}$ leads to cooperation.



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(a) Drone swarm

(b) Path routing



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Action Constrained Multi-Agent Multi-Armed Bandits (1/2)

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$$k_*^{(m)} := \arg \max_{k \in \mathcal{K}^{(m)}} \mu(k)$$

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T Rounds: in each round $t \leq T$

Each agent *m* pulls an arm $k_t^{(m)} \in \mathcal{K}^{(m)}$ and collects reward $X_k^{(m)}(k_t^{(m)})$.

Action Constrained Multi-Agent Multi-Armed Bandits (2/2)

Group regret with K arms M Agents T Rounds

$$\mathbb{E}[\mathsf{R}_{\mathsf{T}}(\mathcal{A})] \coloneqq \sum_{m \in \mathcal{M}} \underbrace{\sum_{t \in \mathcal{T}} (\mu(k_*^{(m)}) - \mu(k_t^{(m)}))}_{\text{agent } m \text{'s regret}}$$

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$$= \sum_{m \in \mathcal{M}} \sum_{t \in \mathcal{T}} \Delta^{(m)}(k)$$

Δ^(m)(k) := μ(k^(m)_{*}) − μ(k) is the reward gap of arm k with respect to agent m's local optimal arm k^(m)_{*}.

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$$= \sum_{m \in \mathcal{M}} \sum_{t \in \mathcal{T}} \Delta^{(m)}(k)$$
$$= \sum_{m \in \mathcal{M}} \sum_{k \in \mathcal{K}^{(m)}} \Delta^{(m)}(k) n_{\mathcal{T}}^{(m)}(k)$$

• $\Delta^{(m)}(k) := \mu(k_*^{(m)}) - \mu(k)$ is the reward gap of arm k with respect to agent m's local optimal arm $k_*^{(m)}$.

• $n_T^{(m)}(k)$ is the number of times that agent *m* pulls arm *k* till the end.

Individual v.s. Cooperative: Algorithm Design

1 Individual UCB arm pull policy: at time *t*, agent *m* pulls



where $n_t^{(m)}(k)$ is the number of times that agent *m* pulls arm *k* up to time *t*, and $\hat{\mu}_t^{(m)}(k)$ is the average of these $n_t^{(m)}(k)$'s observations.

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1 Individual UCB arm pull policy: at time *t*, agent *m* pulls

$$k_t^{(m)} = \underset{k \in \mathcal{K}^{(m)}}{\arg \max} \operatorname{UCB}_t^{(m)}(k) = \underset{k \in \mathcal{K}^{(m)}}{\arg \max} \underbrace{\hat{\mu}_t^{(m)}(k)}_{\text{empirical mean}} + \underbrace{\sqrt{\frac{2 \log t}{n_t^{(m)}(k)}}}_{\text{confidence radius}},$$

where n_t^(m)(k) is the number of times that agent *m* pulls arm *k* up to time *t*, and
µ_t^(m)(k) is the average of these n_t^(m)(k)'s observations.
2 Cooperative UCB arm pull policy: at time *t*, agent *m* pulls

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where $n_t(k)$ is the number of times that all *M* agents pull arm *k* up to time *t*, and $\hat{\mu}_t(k)$ is the average of these $n_t(k)$'s observations.

Individual v.s. Cooperative: Regret Analysis

1 Individual UCB's regret:

$$\mathbb{E}[\mathsf{R}_{\mathsf{T}}(\mathcal{A})] \leqslant O\left(\sum_{m \in \mathcal{M}} \sum_{k \in \mathcal{K}^{(m)} \setminus \{k_*^{(m)}\}} \frac{\log T}{\Delta^{(m)}(k)}\right),$$

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2 Cooperative UCB's regret:

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where $\tilde{\Delta}(k) := \min_{m:k \in \mathcal{K}^{(m)}} \Delta^{(m)}(k)$ is the smallest reward gap of local suboptimal arm *k* with respect to any feasible agent *m*'s local optimal arm.

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2 Cooperative UCB's regret: How good?

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$$\mathbb{E}[\mathsf{R}_{\mathsf{T}}(\mathcal{A})] = \sum_{m \in \mathcal{M}} \underbrace{\sum_{k \in \mathcal{K}^{(m)}} \Delta^{(m)}(k) n_{\mathsf{T}}^{(m)}(k)}_{\text{rearet of agent } m}$$

Regret Lower Bound and Free Exploration Intuition (1/2) $\mathbb{E}[\mathsf{R}_{\mathsf{T}}(\mathcal{A})] = \sum_{m \in \mathcal{M}} \underbrace{\sum_{k \in \mathcal{K}^{(m)}} \Delta^{(m)}(k) n_{\mathsf{T}}^{(m)}(k)}_{\text{regret of agent } m}$

In single agent *m*'s bandit, to distinguish a suboptimal arm requires $n_T^{(m)}(k) = \Omega\left(\frac{\log T}{(\Delta^{(m)}(k))^2}\right) \text{ pulls.}$

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 In single agent *m*'s bandit, to distinguish a suboptimal arm requires n^(m)_T(k) = Ω (log T / (Δ^(m)(k))²) pulls.
 Therefore, Ω (Σ_{k∈K^(m)} Δ^(m)(k) × log T / (Δ^(m)(k))²) = Ω (Σ_{k∈K^(m)} log T / Δ^(m)(k)) regret lower bound.
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 In cooperative multi-agent bandits, we need Ω (log T / (Δ̃(k))²) pulls.
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In single agent m's bandit, to distinguish a suboptimal arm requires $n_T^{(m)}(k) = \Omega\left(\frac{\log T}{(\Lambda^{(m)}(k))^2}\right)$ pulls. • Therefore, $\Omega\left(\sum_{k \in K^{(m)}} \Delta^{(m)}(k) \times \frac{\log T}{(\Delta^{(m)}(k))^2}\right) = \Omega\left(\sum_{k \in K^{(m)}} \frac{\log T}{\Delta^{(m)}(k)}\right)$ regret lower bound. In cooperative multi-agent bandits, we need $\Omega\left(\frac{\log I}{(\tilde{\Delta}(k))^2}\right)$ pulls. 2 • Therefore, $\Omega\left(\sum \tilde{\Delta}(k) \times \frac{\log T}{(\tilde{\Delta}(k))^2}\right) = \Omega\left(\sum_{i \in T} \frac{\log T}{\tilde{\Delta}(k)}\right)$ regret lower bound?

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 Denote F := {k_{*}^(m) : m ∈ M} as arms can be freely explored.

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4 Cooperative UCB is not optimal.

• $\mathcal{K} \setminus \mathcal{F} \subseteq \bigcup_{m \in \mathcal{M}} (\mathcal{K}^{(m)} \setminus \{k_*^{(m)}\})$ of Cooperative UCB's upper bound.

	μ (1)	μ (2)	μ (3)
Agent 1	 	 	
Agent 2	×	>	>
Agent 3	×	×	<

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	Agent 1	 	 	 	
	Agent 2	×	 	>	
	Agent 3	×	×	>	
UCB [Auer, 2002]					
CO-UCB [Yang et al., 2022]					
FreeExp (our work)					

		μ (1)	μ (2)	μ (3)	
	Agent 1	 	 	 	
	Agent 2	×	>	 Image: A start of the start of	
	Agent 3	×	×	 	
UCB [Auer, 2002]		0(($\left(\frac{1}{\Delta(1,2)}\right)$	$+\frac{1}{\Delta(1,3)}$	$\left(1 + \frac{1}{\Delta(2,3)}\right)\log T$
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FreeExp (our work)		<i>O</i> (1))		



1 $\mathcal{K} = \mathcal{F}$, all arms can be freely explored: FreeExp achieves constant regret.

Algorithm 1 The FreeExp Algorithm (for Agent *m*)

- 1: for each time slot t do
- 2: $I_t^{(m)} \leftarrow \arg\max_{k \in \mathcal{K}^{(m)}} \hat{\mu}_t(k)$ \triangleright identify empirical optimal arm
- 3: Send $I_t^{(m)}$ to other agents and collect their $I_t^{(m')}$

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4: $\mathcal{D}_t^{(m)} \leftarrow \{k \in \mathcal{K}^{(m)} : \cup_{CB_t}(k) > \hat{\mu}_t(I_t^{(m)})\}$ \triangleright choose high KL-UCB arms
5: $\mathcal{D}_t^{(m)} \leftarrow \mathcal{D}_t^{(m)} \setminus \{I_t^{(m')} : \forall m' \in \mathcal{M}\}$ \triangleright free exploration
6: if $\mathcal{D}_t^{(m)} = \emptyset$ then
7: $J_t^{(m)} \leftarrow I_t^{(m)} \triangleright$ exploit, if correct only agent m pulls $k_*^{(m)}$
8: else
9: $J_t^{(m)} \leftarrow \begin{cases} I_t^{(m)} & w.p. \frac{1}{2} \\ winformly pick an arm from $\mathcal{D}_t^{(m)} & w.p. \frac{1}{2} \end{cases}$ \triangleright explore$

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Theorem (FreeExp's Regret Upper Bound)

$$\mathbb{E}[\mathsf{R}_{\mathsf{T}}(\mathcal{A})] \leqslant O\left(\sum_{k \in \mathcal{K} \setminus \mathcal{F}} \frac{\log T}{\tilde{\Delta}(k)}\right) + O\left(\sum_{k \in \mathcal{F}} 1\right)$$

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1 Theoretical improvement: summation range

$$\underbrace{(m \times k) \in (\mathcal{M} \times \mathcal{K})}_{\text{UCB [Auer, 2002]}} \Longrightarrow \underbrace{k \in \bigcup_{m \in \mathcal{M}} (\mathcal{K}^{(m)} \setminus \{k_*^{(m)}\})}_{\text{CO-UCB [Yang et al., 2022]}} \Longrightarrow \underbrace{k \in \mathcal{K} \setminus \mathcal{F}}_{\text{FreeExp (ours)}}$$

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2 Regret optimality: Match regret lower bound up to constant coefficients.

$$\mathbb{E}[\mathsf{R}_{\mathsf{T}}(\mathcal{A})] \geqslant \Omega\left(\sum_{k \in \mathcal{K} \setminus \mathcal{F}} rac{\log T}{\widetilde{\Delta}(k)}
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Similar Finite regret in special case: When $\mathcal{K} = \mathcal{F}$ (all arms are free), the regret reduces O(1).

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1 Regret due to free arms in \mathcal{F}

2 Regret due to non-free arms in $\mathcal{K} \setminus \mathcal{F}$

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1 Regret due to free arms in \mathcal{F}

Finite # time slots
$$\{l_t^{(m)}: m \in \mathcal{M}\} \neq \mathcal{F} = \{k_*^{(m)}: m \in \mathcal{M}\} \Longrightarrow$$
 Finite regret

estimated free arm set

true free arm set

2 Regret due to non-free arms in $\mathcal{K} \setminus \mathcal{F}$

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estimated free arm set **2** Regret due to non-free arms in $\mathcal{K} \setminus \mathcal{F}$ $\sum_{k \in \mathcal{K} \setminus \mathcal{F}} \sum_{m \in \mathcal{M}} \Delta^{(m)}(k) n_T^{(m)}(k) \text{ and } \sum_{m \in \mathcal{M}} n_T^{(m)}(k) \leqslant \frac{\log T}{(\tilde{\Delta}(k))^2} \implies \frac{\log T}{\tilde{\Delta}(k)} \text{ regret}$

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estimated free arm set **Regret due to non-free arms in** $\mathcal{K} \setminus \mathcal{F}$

$$= \sum_{k \in \mathcal{K} \setminus \mathcal{F}} \sum_{m \in \mathcal{M}} \Delta^{(m)}(k) n_T^{(m)}(k) \text{ and } \sum_{m \in \mathcal{M}} n_T^{(m)}(k) \leqslant \frac{\log T}{(\tilde{\Delta}(k))^2} \Longrightarrow \frac{\log T}{\tilde{\Delta}(k)} \text{ regret }$$

• $\sigma(\ell)$: sort agents in a descending order based on the magnitude of $\Delta^{(m)}(k)$

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1 Regret due to free arms in \mathcal{F}

Finite # time slots
$$\{I_t^{(m)}: m \in \mathcal{M}\} \neq \mathcal{F} = \{k_*^{(m)}: m \in \mathcal{M}\} \Longrightarrow$$
 Finite regre

true free arm set

estimated free arm set **2** Regret due to non-free arms in $\mathcal{K} \setminus \mathcal{F}$

$$= \sum_{k \in \mathcal{K} \setminus \mathcal{F}} \sum_{m \in \mathcal{M}} \Delta^{(m)}(k) n_T^{(m)}(k) \text{ and } \sum_{m \in \mathcal{M}} n_T^{(m)}(k) \leqslant \frac{\log T}{(\tilde{\Delta}(k))^2} \Longrightarrow \frac{\log T}{\tilde{\Delta}(k)} \text{ regret }$$

• $\sigma(\ell)$: sort agents in a descending order based on the magnitude of $\Delta^{(m)}(k)$ • $\sum_{\ell=1}^{L} n_{T}^{(\sigma(\ell))}(k) \leq \frac{\log T}{(\Delta^{(\sigma(L))}(k))^{2}} \xrightarrow{\text{Abel's summation}} \sum_{\ell=1}^{L} \Delta^{(\sigma(\ell))}(k) n_{T}^{(\sigma(\ell))}(k) \leq \frac{C \log T}{\Delta^{(\sigma(L))}(k)}$

Simulations (1/2): FreeExp vs. Baselines



Figure 2: FreeExp vs. baselines

Although with tighter theoretical performance, the empirical performance of FreeExp is not as good as CO-KL-UCB.

Simulations (2/2): Vary parameters of MA2B-HR



Figure 3: Vary parameters of MA2B-HR

■ In Figure (c), the more free arms, the better regret of FreeExp.

Conclusion

- 1 Discover the **free exploration mechanism** in multi-agent bandits with action constraints model.
- 2 Propose a new **regret lower bound**, echoing the free exploration mechanism.
- **3** Devise the **FreeExp** algorithm utilizing the free exploration mechanism.
- **4** Prove that FreeExp's regret upper bound tightly matches the lower bound.
- **5** Conduct **simulations** to validate **FreeExp**'s empirical performance.

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Future works:

- Fairness among heterogeneous agents?
- **Reduce communications from** O(T) to $O(\log T)$?

Thank you!

Full paper at openreview.net/pdf?id=8kKEz1bnIEp

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